Effective Computation of Hodge Cycles

Edgar Costa (MIT) August 27, 2024, Simons Symposium on Geometry of Non-Closed Fields

Joint work with Nicholas Mascot, Jeroen Sijsling, John Voight, and Emre Can Sertöz

Endomorphism ring of an abelian variety

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Given *A* compute the endomorphism ring End *A*.

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• Over a finite field, Honda–Tate theory tells us

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From the equations of *A* determine a basis for End *A* and their equations in *A*×*A*.

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determines the *k*-isogeny class and the isomorphism class of $\text{End}(A) \otimes \mathbb{Q}$.

- There are several in principle algorithms to do this over a number field. These involve, a day/night algorithm:
	- by day: search for reasonable morphisms;
	- by night: restrict your search space.

Our setup

Let *C* be a nice (smooth, projective, geometrically integral) curve over *k* of genus *g* given by equations. Let *J* be the Jacobian of *C*.

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But why?

- It is an interesting challenge [*citation needed*].
- If End *J* contains non-trivial idempotents, we can hope to decompose *J* into abelian varieties of smaller dimension.
- If End *J* is non-trivial, then this allows us to find a modular form that describes the arithmetic properties of *J* and *C*.
- An algorithm to decide transcendence of 1-periods using Huber–Wüstholz theory (Ouaknine–Worrell–Sertöz)

An analytic description of the Jacobian

Via a chosen embedding of *k* into C, we can consider *C* as a Riemann surface, and

 $J_{\mathbb{C}} = H^0(C, \Omega_C)^{\vee}/H_1(C, \mathbb{Z}) = \mathbb{C}^g/\Lambda,$

 ω where we pick an k basis for $H^0(\mathcal{C},\Omega_{\mathcal{C}})=k\omega_1\oplus\ldots\oplus k\omega_g,$ hence,

$$
\Lambda = \left\{ \left(\int_{\gamma} \omega_1, \ldots, \int_{\gamma} \omega_g \right) \in \mathbb{C}^g \ : \ \gamma \in H_1(C, \mathbb{Z}) \right\} \cong \mathbb{Z}^{2g}.
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In other words, *J* is a complex torus (plus a polarization).

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- Using Λ, we can hope to understand *J* analytically…

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In other words, *J* is a complex torus (plus a polarization).

- We can calculate Λ numerically.
- Using Λ, we can hope to understand *J* analytically…
- and perhaps even to be able to transfer these results to the algebraic setting.

By picking a *k*-basis for $H^0(C, \Omega_C)$, we have

 $\text{End}(J) = \{T \in M_q(k) \mid T\Lambda \subset \Lambda\}$

Hence, if Π is a period matrix for *C*, i.e., $\Lambda = \Pi \mathbb{Z}^{2g}$, then we are reduced to finding a Z-basis of the solutions (*T*, *R*) to

$$
T\Pi = \Pi R, \qquad T \in M_g(k^{\text{al}}), \quad R \in M_{2g}(\mathbb{Z}).
$$

Heuristically, via lattice reduction algorithms, we can find such a \mathbb{Z} -basis.

There is no obvious way to prove that our guesses are actually correct...

Representing endomorphisms

$$
\alpha_C : C \xrightarrow{A} J \xrightarrow{\alpha} J \longrightarrow \text{Sym}^g(C)
$$

$$
P \longmapsto \{Q_1, \dots, Q_g\} \iff \alpha([P - P_0]) = \left[\sum_{i=1}^g Q_i - P_0\right]
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This traces out a divisor on $C \times C$, which determines α .

Given $\alpha\in\mathrm{M}_g(k^{\mathsf{al}})$ this divisor is a certificate of containment $[\alpha\bullet]$ for $\alpha\in\mathsf{End}\,J^{\mathsf{al}}.$

Theorem (C–Mascot–Sijsling–Voight)

We give an algorithm for

$$
M_g(k^{al}) \ni \alpha \mapsto \begin{cases} \text{true} & \text{if } \alpha \in \text{End } J^{al}, \text{ and a certificate} \\ \text{false} & \text{if } \alpha \notin \text{End } J^{al} \end{cases}
$$

By interpolation via α_c or by locally solving a differential equation on $C \times C$.

We give an algorithm that computes <code>End</code> J $^{\mathsf{al}}$ *with a certificate* $\boxed{\checkmark}$ *.*

This is a day/night algorithm:

• By day, we compute Λ ⊂ C *^g* numerically and then certify *B* ⊆ End *J* al .

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This is a day/night algorithm:

- \cdot By day, we compute $\Lambda\subset\mathbb{C}^g$ numerically and then certify $B\subseteq \mathsf{End}\, J^{\sf al}.$
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 $\{L_{p_1}(t) := \det(1 - t \operatorname{Frob}_p | H^1), L_{p_2}(t), \ldots, L_{p_i}(t)\} \longmapsto$ upper bounds on End J^{al}

- \cdot The $L_p(t)$ polynomials are as random as End *J*^{al} allows it.
- Two polynomials *Lp*(*t*) and *Lq*(*t*) suffice to obtain a sharp upperbound.
- For (*p*, *q*) in a set of positive density, but unknown apriori.

We give an algorithm that computes <code>End</code> J $^{\mathsf{al}}$ *with a certificate* $\boxed{\checkmark}$ *.*

This is a day/night algorithm:

- \cdot By day, we compute $\Lambda\subset\mathbb{C}^g$ numerically and then certify *B* ⊆ **End** J^{al}.
- By night, we search for evidence that End *J* al ⊆ *B*.

 $Frob_p \mod p^N \; \bigodot_h H^1_{\text{crys}}(\mathcal{C},\mathbb{Z}_p) \longmapsto \text{upper bounds on } \mathsf{End}\, J^{\text{al}}$

- ∙ Frob_{*p*} mod p^N is a byproduct of computing $L_p(t) = det(1 t \operatorname{Frob}_p|H_{MW}^1).$
- We check what correspondences $C \rightsquigarrow C$ mod p lift to $C \rightsquigarrow C$ mod p^N .

Examples

- Our method works just as well for isogenies and projections.
- \cdot We have verified, decomposed and matched the 66, 158 curves over $\mathbb O$ of genus 2 in the *L-functions and modular form database* (LMFDB).
- \cdot The algorithms verify that the plane quartic

$$
C: x4 - x3y + 2x3z + 2x2yz + 2x2z2 - 2xy2z + 4xyz2 - y3z + 3y2z2 + 2yz3 + z4 = 0
$$

has complex multiplication.

• Try it:

<https://github.com/edgarcosta/endomorphisms>

contains friendly button-push algorithms.

Picard lattice of a K3 surface

Let *X* be a K3 surface defined over $k \subset \mathbb{C}$. We view *X* also as a complex manifold. $Pic X^{al} \simeq \mathbb{Z}\langle \text{algebraic curves in } X \rangle / \langle \text{linear equivalences} \rangle \subset H_2(X, \mathbb{Z})$

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Let *X* be a K3 surface defined over *k* ⊂ C. We view *X* also as a complex manifold. $Pic X^{al} \simeq \mathbb{Z}\langle \text{algebraic curves in } X \rangle / \langle \text{linear equivalences} \rangle \subset H_2(X, \mathbb{Z})$

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From the equations of *X*, compute Pic $X^{\mathsf{al}} \subset H_2(X, \mathbb{Z})$ as a $\mathsf{Gal}(k^{\mathsf{al}}/k)$ -module.

- \cdot Over finite field, Tate conjecture tells us that $\mathsf{det}(1-t\,\mathsf{Frob}|H^2(X,\mathbb{Q}_\ell))\in \mathbb{Z}[t]$ gives us the rank of Pic *X*.
- There are several in principle algorithms to compute rk Pic *X* or even Pic *X* over a number field.

These involve, a day/night algorithm:

- by day: find curve classes in Pic *X*;
- \cdot by night: restrict the ambient space for Pic *X* ⊂ $H^2(X,\mathbb{Z})$.

"The evaluation of ρ *for a given surface presents in general grave difficulties."* — Zariski

An analytic approach

Lefschetz (1,1) theorem

A homology class $\gamma\in H_2(X,{\mathbb Z})$ is in Pic $X^{\sf al}$ if and only if $\int_\gamma\omega_X=$ 0, where ω_X is the nonzero holomorphic 2-form ω_X on *X*, unique up to scaling.

Hence, if **Π** is the period vector for ω_X , i.e., $[\int_\gamma \omega_\mathsf{X}]_{\gamma\in H_2(\mathsf{X},\mathbb{Z})}$, then we are reduced to finding a lattice $\Lambda \subset H_2(X, \mathbb{Z})$ of solutions

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\Pi R = 0, \qquad R \in \mathbb{Z}^{22}.
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- Heuristically, via lattice reduction algorithms, we can find Λ.
- There is no obvious way to prove that our guesses are actually correct...
- Nonetheless, a posteriori, one can compute $B \gg 0$ such that

 $Pic(X^{al})_{|B} := \mathbb{Z}\langle \gamma \in Pic(X^{al}) \mid -\gamma_{prim}^2 < B \rangle \subseteq \Lambda$ (Lairez–Sertöz).

$$
X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3
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• It is a fiber in a pencil that has generic rank 19 and matching upper bounds can be deduced by positive characteristic methods.

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- \cdot Heuristically, one computes Λ such that $\mathsf{Pic}(X^{\mathsf{al}})_{|B} \subseteq \Lambda \overset{?}{\subseteq} \mathsf{Pic} X^{\mathsf{al}}.$
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- No small rational curves: There are no lines, no conics, no twisted cubics.
- The "smallest" non-trivial curves that appear are smooth rational quartics.
- Lattice computations with Λ predict that there are

133056

smooth rational quartics spanning Λ.

Reconstructing isolated curves from their Hodge classes

Turns out one can compute a bit more for hypersurfaces

$$
\varphi: H_2(X,\mathbb{Z}) \times H^2_{\mathrm{dR}}(X/k) \to \mathbb{C} \qquad (\gamma,\omega) \mapsto \int_{\gamma} \omega
$$

Note, if $\gamma \in \text{Pic } X^{\text{al}}$, then $\frac{1}{2\pi i} \int_{\gamma} \omega \in k^{\text{al}}$ for $\omega \in F^1H^2_{\text{dR}}(X/k)$.

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Theorem (Movasati–Sertöz)

If $\gamma = [Y] \in H_2(X, \mathbb{Z})$ for a curve $Y \subset X$ then from $\frac{1}{2\pi i} (\int_{\gamma} \omega)_{\omega \in F^1}$ one can construct an ideal I_{γ} such that $I(Y) \subseteq I_{\gamma}$.

In favorable circumstances we expect low order equations in I_{γ} to span *I*(*Y*).

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Theorem (Cifani–Pirola–Schlesinger)

For a smooth rational quartic *Y* \subset *X* we have that the equation of the quadric containing *Y* generates $I_{[Y],2}$, i.e., $I(Y)_2 = I_{[Y],2}$.

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Reconstruct the quadrics containing some of the 133056 smooth rational quartics in *X* using the curve classes.

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• Fortunately, there is a small Aut(Λ) orbit of size 336.

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Goal

Reconstruct the quadrics containing some of the 133056 smooth rational quartics in *X* using the curve classes.

- Fortunately, there is a small Aut(Λ) orbit of size 336.
- Hence, we expect an orbit of 168 quadrics each containing a pair of quartics.
- We aim reconstruct the ten (algebraic!) coefficients of these quadrics.

Goal

Reconstruct the ten coefficients of these quadrics in a Galois orbit of size 168.

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- The minimal polynomials of these elements have incredibly large height.

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− 1268317331496745879603035032448157273146519836562713924560050631153969519297207668270922371313*x* ¹⁴⁷ + · · ·

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- Every computation must be done very selectively.
- We solve the isomorphism problem between the different presentations by refining the complex embeddings and inverting a Vandermonde matrix. The abstract isomorphism problem feels hopeless otherwise.

Intersecting the quadric with *X*

$$
Q: a_0x^2 + a_1xy + \cdots + a_9w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168
$$

Goal

Show that *Q* ∩ *X* decomposes into two quartic curves.

• It suffices to show that the singular locus *S* of *Q* ∩ *X* consists of 10 distinct reduced points.

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- Hopeless to do this directly!
- Working over \mathbb{F}_p we find 10 distinct points. Hence, *S* is zero-dimensional and reduced, and $\deg S \leq 10$.

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- Hopeless to do this directly!
- Working over F*^p* we find 10 distinct points. Hence, *S* is zero-dimensional and reduced, and $\deg S \leq 10$.
- \cdot We conclude deg $S = 10$ via Gotzmann regularity theorem, by checking that dim $L[x, y, z, w]_{\bullet}/I_{\bullet} = 10$ for $\bullet = 6, 7$, where $V(I) = S$.

$$
\Lambda_Q:=\langle [C]:C\subset \sigma(Q)\cap X,\,\sigma:L\hookrightarrow \mathbb{C}\rangle\subseteq \text{Pic}(X^{\text{al}})_{|B}\subseteq \Lambda\stackrel{?}{\subseteq} \text{Pic} X^{\text{al}}
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The inclusion $\Lambda_{\Omega} \subseteq \Lambda$ is not explicit.

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The inclusion Λ ^{\circ} \subset Λ is not explicit.

Nonetheless, Pic *X* al and Λ are saturated in *H*2(*X*, Z).

Hence, it is sufficient to show that $rk \Lambda_0 = rk \Lambda = 19$.

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We can do this in two ways:

 \cdot Compute the intersections of these 336 curves with each other over \mathbb{F}_p .

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Hence, it is sufficient to show that $rk \Lambda_{\Omega} = rk \Lambda = 19$.

We can do this in two ways:

- \cdot Compute the intersections of these 336 curves with each other over \mathbb{F}_p .
- Certify that these correspond to the original classes. Showing that there are at most 66528 distinct quadrics. Can be done over C. This establishes a bijection between these quadrics and the 168 pairs of quartic curve classes that they correspond to.

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Goal

Compute *K* and Gal(*K*/Q) acting on Λ*Q*.

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Via the identification with the original classes we have $\frac{1}{2\pi i}\left(\int_C\omega\right)_{\omega\in F^1}\in K^{21}.$

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Via the identification with the original classes we have $\frac{1}{2\pi i}\left(\int_C\omega\right)_{\omega\in F^1}\in K^{21}.$ These can be reconstructed in the same fashion as we reconstructed *aⁱ* . Unclear how to certify such heuristic guesses!

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Q: a_0x^2 + a_1xy + \cdots + a_9w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168
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Even if given the order $O \subset K$ over which the quartics are defined over, no obvious control over denominators of $\frac{1}{2\pi i}\int_C \omega$.

Can one compute *K* using geometry without Gröbner basis?

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x^{24} + x^{22} - 24x^{21} - 84x^{20} - 205x^{19} - 155x^{18} - 770x^{17} - 500x^{16} + 18916x^{15} + 36988x^{14} + 109234x^{13} + 387901x^{12} + 373961x^{11} - 18170x^{10} + 75132x^9 + 10381x^8 - 123071x^7 + 108274x^6 - 41580x^5 + 39936x^4 - 21911x^3 + 4032x^2 + 1428x + 616
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and $Gal(F/\mathbb{Q}) = C_3 \times PGL(2,7)$ (with size 14 times smaller than Aut Pic X^{al}). Do we have Gal(K/\mathbb{Q}) $\stackrel{?}{=}$ Aut Λ? Can we compute Gal(K/\mathbb{Q}) by hand?

Theorem (C–Sertöz)

The quartic surface $X : x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$ has $\mathsf{Pic}\,X^{\mathsf{al}} = \mathsf{\Lambda},$ generated by quartics over a quadratic extension of $L:=\mathbb{Q}(\{a_i\}_i).$

We are hoping to streamline this method and also figure out its applications. Hopefully, also be able handle families, e.g.,

$$
X: x^4 + txyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3(\mathbb{Q}(t))
$$

Do you have a challenge K3 surface for us?