Effective Computation of Hodge Cycles

Edgar Costa (MIT) August 27, 2024, Simons Symposium on Geometry of Non-Closed Fields

Joint work with Nicholas Mascot, Jeroen Sijsling, John Voight, and Emre Can Sertöz

Endomorphism ring of an abelian variety

Let A be an abelian variety defined over k.

Goal

Given A compute the endomorphism ring End A.

Endomorphism ring of an abelian variety

Let A be an abelian variety defined over k.

Goal

Given A compute the endomorphism ring End A.

• Over a finite field, Honda–Tate theory tells us

 $\det(1-t\operatorname{Frob}|H^1(A,\mathbb{Q}_\ell)\in\mathbb{Z}[t]$

determines the *k*-isogeny class and the isomorphism class of $End(A) \otimes \mathbb{Q}$.

Endomorphism ring of an abelian variety

Let A be an abelian variety defined over k.

Goal

From the equations of A determine a basis for End A and their equations in $A \times A$.

• Over a finite field, Honda–Tate theory tells us

 $\det(1 - t \operatorname{Frob}|H^1(A, \mathbb{Q}_\ell) \in \mathbb{Z}[t]$

determines the *k*-isogeny class and the isomorphism class of $End(A) \otimes \mathbb{Q}$.

- There are several in principle algorithms to do this over a number field. These involve, a day/night algorithm:
 - by day: search for reasonable morphisms;
 - by night: restrict your search space.

Our setup

Let *C* be a nice (smooth, projective, geometrically integral) curve over *k* of genus *g* given by equations. Let *J* be the Jacobian of *C*.

Goal

Given the equations of C, compute the endomorphism ring End J^{al}.

Our setup

Let *C* be a nice (smooth, projective, geometrically integral) curve over *k* of genus *g* given by equations. Let *J* be the Jacobian of *C*.

Goal

Given the equations of C, compute the endomorphism ring $\operatorname{End} J^{\operatorname{al}}$.

But why?

- It is an interesting challenge [*citation needed*].
- If End J contains non-trivial idempotents, we can hope to decompose J into abelian varieties of smaller dimension.
- If End J is non-trivial, then this allows us to find a modular form that describes the arithmetic properties of J and C.
- An algorithm to decide transcendence of 1-periods using Huber–Wüstholz theory (Ouaknine–Worrell–Sertöz)

An analytic description of the Jacobian

Via a chosen embedding of k into \mathbb{C} , we can consider C as a Riemann surface, and

 $J_{\mathbb{C}} = H^0(C, \Omega_C)^{\vee} / H_1(C, \mathbb{Z}) = \mathbb{C}^g / \Lambda,$

where we pick an k basis for $H^0(C, \Omega_C) = k\omega_1 \oplus \ldots \oplus k\omega_g$, hence,

$$\Lambda = \left\{ \left(\int_{\gamma} \omega_1, \dots, \int_{\gamma} \omega_g \right) \in \mathbb{C}^g : \gamma \in H_1(C, \mathbb{Z}) \right\} \cong \mathbb{Z}^{2g}.$$

In other words, *J* is a complex torus (plus a polarization).

- \cdot We can calculate Λ numerically.
- · Using Λ , we can hope to understand J analytically...

An analytic description of the Jacobian

Via a chosen embedding of k into \mathbb{C} , we can consider C as a Riemann surface, and

 $J_{\mathbb{C}} = H^0(C, \Omega_C)^{\vee} / H_1(C, \mathbb{Z}) = \mathbb{C}^g / \Lambda,$

where we pick an k basis for $H^0(C, \Omega_C) = k\omega_1 \oplus \ldots \oplus k\omega_g$, hence,

$$\Lambda = \left\{ \left(\int_{\gamma} \omega_1, \ldots, \int_{\gamma} \omega_g \right) \in \mathbb{C}^g : \gamma \in H_1(C, \mathbb{Z}) \right\} \cong \mathbb{Z}^{2g}.$$

In other words, *J* is a complex torus (plus a polarization).

- \cdot We can calculate Λ numerically.
- · Using Λ , we can hope to understand J analytically...
- and perhaps even to be able to transfer these results to the algebraic setting.

By picking a *k*-basis for $H^0(C, \Omega_C)$, we have

 $\operatorname{End}(J) = \{T \in M_g(k) \mid T\Lambda \subset \Lambda\}$

Hence, if Π is a period matrix for *C*, i.e., $\Lambda = \Pi \mathbb{Z}^{2g}$, then we are reduced to finding a \mathbb{Z} -basis of the solutions (T, R) to

$$T\Pi = \Pi R, \qquad T \in M_g(k^{al}), \quad R \in M_{2g}(\mathbb{Z}).$$

Heuristically, via lattice reduction algorithms, we can find such a \mathbb{Z} -basis. There is no obvious way to prove that our guesses are actually correct...

Representing endomorphisms

$$\alpha_{\mathcal{C}} : \mathcal{C} \xrightarrow{AJ} J \xrightarrow{\alpha} J \dashrightarrow \mathsf{Sym}^{g}(\mathcal{C})$$
$$P \longmapsto \{Q_{1}, \dots, Q_{g}\} \Longleftrightarrow \alpha([P - P_{0}]) = \left[\sum_{i=1}^{g} Q_{i} - P_{0}\right]$$

This traces out a divisor on $C \times C$, which determines α .

Representing endomorphisms

$$\alpha_{\mathcal{C}} : \mathcal{C} \xrightarrow{AJ} J \xrightarrow{\alpha} J \dashrightarrow \operatorname{Sym}^{g}(\mathcal{C})$$
$$P \longmapsto \{Q_{1}, \dots, Q_{g}\} \Longleftrightarrow \alpha([P - P_{0}]) = \left[\sum_{i=1}^{g} Q_{i} - P_{0}\right]$$

This traces out a divisor on $C \times C$, which determines α .

Given $\alpha \in M_g(k^{al})$ this divisor is a certificate of containment $\alpha \in \text{End } J^{al}$. **Theorem (C-Mascot-Sijsling-Voight)** We give an algorithm for $M_g(k^{al}) \ni \alpha \mapsto \begin{cases} true & \text{if } \alpha \in \text{End } J^{al}, \text{and a certificate } \alpha \in false & \text{if } \alpha \notin \text{End } J^{al} \end{cases}$

By interpolation via α_{C} or by locally solving a differential equation on $C \times C$.

We give an algorithm that computes $\operatorname{End} J^{\operatorname{al}}$ with a certificate $\boxed{\checkmark_{\bullet}}$.

This is a day/night algorithm:

• By day, we compute $\Lambda \subset \mathbb{C}^g$ numerically and then certify $B \subseteq \operatorname{End} J^{\operatorname{al}}$.

We give an algorithm that computes $\operatorname{End} J^{\operatorname{al}}$ with a certificate $\boxed{\checkmark_{\bullet}}$.

This is a day/night algorithm:

- By day, we compute $\Lambda \subset \mathbb{C}^g$ numerically and then certify $B \subseteq \operatorname{End} J^{\operatorname{al}}$.
- By night, we search for evidence that $\operatorname{End} J^{\operatorname{al}} \subseteq B$.

We give an algorithm that computes $\operatorname{End} J^{\operatorname{al}}$ with a certificate $\boxed{\checkmark_{\bullet}}$.

This is a day/night algorithm:

- By day, we compute $\Lambda \subset \mathbb{C}^g$ numerically and then certify $B \subseteq \operatorname{End} J^{\operatorname{al}}$.
- By night, we search for evidence that $\operatorname{End} J^{\operatorname{al}} \subseteq B$.

We give an algorithm that computes End J^{al} with a certificate \bigvee_{\blacksquare} .

This is a day/night algorithm:

• By day, we compute $\Lambda \subset \mathbb{C}^g$ numerically and then certify $B \subseteq \operatorname{End} J^{\operatorname{al}}$.

$$M_{g}(k^{\mathsf{al}}) \ni \alpha \mapsto \begin{cases} \mathsf{true} & \text{if } \alpha \in \mathsf{End}\,J^{\mathsf{al}}, \text{and a certificate } \alpha_{\bullet} \\ \mathsf{false} & \text{if } \alpha \notin \mathsf{End}\,J^{\mathsf{al}} \end{cases}$$

• By night, we search for evidence that $\operatorname{End} J^{\operatorname{al}} \subseteq B$.

We give an algorithm that computes End J^{al} with a certificate √.

This is a day/night algorithm:

- By day, we compute $\Lambda \subset \mathbb{C}^g$ numerically and then certify $B \subseteq \operatorname{End} J^{\operatorname{al}}$.
- By night, we search for evidence that $\operatorname{End} J^{\operatorname{al}} \subseteq B$.

 $\{L_{p_1}(t) := \det(1 - t \operatorname{Frob}_p | H^1), L_{p_2}(t), \dots, L_{p_i}(t)\} \longrightarrow \text{upper bounds on } \operatorname{End} J^{\operatorname{al}}$

- The $L_p(t)$ polynomials are as random as End J^{al} allows it.
- Two polynomials $L_p(t)$ and $L_q(t)$ suffice to obtain a sharp upperbound.
- For (p,q) in a set of positive density, but unknown apriori.

We give an algorithm that computes End J^{al} with a certificate <u></u>.

This is a day/night algorithm:

- By day, we compute $\Lambda \subset \mathbb{C}^g$ numerically and then certify $B \subseteq \operatorname{End} J^{\operatorname{al}}$.
- By night, we search for evidence that $\operatorname{End} J^{\operatorname{al}} \subseteq B$.

 $\operatorname{Frob}_{p} \operatorname{mod} p^{N} \, \bigcirc_{\mathcal{H}} \, H^{1}_{\operatorname{crys}}(C, \mathbb{Z}_{p}) \xrightarrow{} \operatorname{upper bounds on } \operatorname{End} J^{\operatorname{al}}$

- Frob_p mod p^N is a byproduct of computing $L_p(t) = \det(1 t \operatorname{Frob}_p | H^1_{MW})$.
- We check what correspondences $C \rightsquigarrow C \mod p$ lift to $C \rightsquigarrow C \mod p^N$.

Examples

- Our method works just as well for isogenies and projections.
- We have verified, decomposed and matched the 66, 158 curves over \mathbb{Q} of genus 2 in the *L*-functions and modular form database (LMFDB).
- \cdot The algorithms verify that the plane quartic

$$C: x^{4} - x^{3}y + 2x^{3}z + 2x^{2}yz + 2x^{2}z^{2} - 2xy^{2}z + 4xyz^{2}$$
$$-y^{3}z + 3y^{2}z^{2} + 2yz^{3} + z^{4} = 0$$

has complex multiplication.

• Try it:

https://github.com/edgarcosta/endomorphisms
contains friendly button-push algorithms.

Picard lattice of a K3 surface

Let X be a K3 surface defined over $k \subset \mathbb{C}$. We view X also as a complex manifold.

 $\operatorname{Pic} X^{\operatorname{al}} \simeq \mathbb{Z} \langle \operatorname{algebraic} \operatorname{curves} \operatorname{in} X \rangle / \langle \operatorname{linear} \operatorname{equivalences} \rangle \subset H_2(X, \mathbb{Z})$

Goal

Given X compute $\operatorname{Pic} X^{\operatorname{al}}$.

Picard lattice of a K3 surface

Let X be a K3 surface defined over $k \subset \mathbb{C}$. We view X also as a complex manifold.

 $\operatorname{Pic} X^{\operatorname{al}} \simeq \mathbb{Z} \langle \operatorname{algebraic} \operatorname{curves} \operatorname{in} X \rangle / \langle \operatorname{linear} \operatorname{equivalences} \rangle \subset H_2(X, \mathbb{Z})$

Goal

Given X compute $Pic X^{al}$.

• Over finite field, Tate conjecture tells us that $det(1 - t \operatorname{Frob}|H^2(X, \mathbb{Q}_{\ell})) \in \mathbb{Z}[t]$ gives us the rank of Pic X.

Let X be a K3 surface defined over $k \subset \mathbb{C}$. We view X also as a complex manifold.

Pic $X^{al} \simeq \mathbb{Z}\langle algebraic \ curves \ in \ X \rangle / \langle linear \ equivalences \rangle \subset H_2(X, \mathbb{Z})$

Goal

From the equations of X, compute $\operatorname{Pic} X^{\operatorname{al}} \subset H_2(X, \mathbb{Z})$ as a $\operatorname{Gal}(k^{\operatorname{al}}/k)$ -module.

- Over finite field, Tate conjecture tells us that $det(1 t \operatorname{Frob}|H^2(X, \mathbb{Q}_{\ell})) \in \mathbb{Z}[t]$ gives us the rank of Pic X.
- There are several in principle algorithms to compute rk Pic X or even Pic X over a number field.

These involve, a day/night algorithm:

- by day: find curve classes in Pic*X*;
- by night: restrict the ambient space for $\operatorname{Pic} X \subset H^2(X, \mathbb{Z})$.

"The evaluation of ho for a given surface presents in general grave difficulties." — Zariski

An analytic approach

Lefschetz (1,1) theorem

A homology class $\gamma \in H_2(X, \mathbb{Z})$ is in Pic X^{al} if and only if $\int_{\gamma} \omega_X = 0$, where ω_X is the nonzero holomorphic 2-form ω_X on X, unique up to scaling.

Hence, if Π is the period vector for ω_X , i.e., $[\int_{\gamma} \omega_X]_{\gamma \in H_2(X,\mathbb{Z})}$, then we are reduced to finding a lattice $\Lambda \subset H_2(X,\mathbb{Z})$ of solutions

$$\Pi R = 0, \qquad R \in \mathbb{Z}^{22}.$$

An analytic approach

Lefschetz (1,1) theorem

A homology class $\gamma \in H_2(X, \mathbb{Z})$ is in Pic X^{al} if and only if $\int_{\gamma} \omega_X = 0$, where ω_X is the nonzero holomorphic 2-form ω_X on X, unique up to scaling.

Hence, if Π is the period vector for ω_X , i.e., $[\int_{\gamma} \omega_X]_{\gamma \in H_2(X,\mathbb{Z})}$, then we are reduced to finding a lattice $\Lambda \subset H_2(X,\mathbb{Z})$ of solutions

$$\Pi R=0, \qquad R\in\mathbb{Z}^{22}.$$

- \cdot Π can be computed via deformation for projective hypersurfaces (Sertöz).
- \cdot Heuristically, via lattice reduction algorithms, we can find $\Lambda.$
- There is no obvious way to prove that our guesses are actually correct...

An analytic approach

Lefschetz (1,1) theorem

A homology class $\gamma \in H_2(X, \mathbb{Z})$ is in Pic X^{al} if and only if $\int_{\gamma} \omega_X = 0$, where ω_X is the nonzero holomorphic 2-form ω_X on X, unique up to scaling.

Hence, if Π is the period vector for ω_X , i.e., $[\int_{\gamma} \omega_X]_{\gamma \in H_2(X,\mathbb{Z})}$, then we are reduced to finding a lattice $\Lambda \subset H_2(X,\mathbb{Z})$ of solutions

$$\Pi R=0, \qquad R\in\mathbb{Z}^{22}.$$

- \cdot Π can be computed via deformation for projective hypersurfaces (Sertöz).
- \cdot Heuristically, via lattice reduction algorithms, we can find $\Lambda.$
- There is no obvious way to prove that our guesses are actually correct...
- Nonetheless, a posteriori, one can compute $B \gg 0$ such that

 $\mathsf{Pic}(X^{\mathsf{al}})_{|B} := \mathbb{Z} \langle \gamma \in \mathsf{Pic}(X^{\mathsf{al}}) \mid -\gamma_{\mathsf{prim}}^2 < B \rangle \subseteq \Lambda \qquad (\mathsf{Lairez-Sert}\"{oz}).$

$$X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$$

• It is a fiber in a pencil that has generic rank 19 and matching upper bounds can be deduced by positive characteristic methods.

$$X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$$

- It is a fiber in a pencil that has generic rank 19 and matching upper bounds can be deduced by positive characteristic methods.
- No known explicit descriptions of Pic X^{al}.

$$X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$$

- It is a fiber in a pencil that has generic rank 19 and matching upper bounds can be deduced by positive characteristic methods.
- No known explicit descriptions of $Pic X^{al}$.
- Heuristically, one computes Λ such that $\operatorname{Pic}(X^{\operatorname{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \operatorname{Pic} X^{\operatorname{al}}$.
- We can compute Aut A, the isomorphism class seems to be $F_7 \times PGL(2,7)$.

$$X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$$

- It is a fiber in a pencil that has generic rank 19 and matching upper bounds can be deduced by positive characteristic methods.
- No known explicit descriptions of $Pic X^{al}$.
- Heuristically, one computes Λ such that $\operatorname{Pic}(X^{\operatorname{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \operatorname{Pic} X^{\operatorname{al}}$.
- We can compute Aut A, the isomorphism class seems to be $F_7 \times PGL(2,7)$.
- No small rational curves: There are no lines, no conics, no twisted cubics.
- The "smallest" non-trivial curves that appear are smooth rational quartics.

$$X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$$

- It is a fiber in a pencil that has generic rank 19 and matching upper bounds can be deduced by positive characteristic methods.
- No known explicit descriptions of $Pic X^{al}$.
- Heuristically, one computes Λ such that $\operatorname{Pic}(X^{\operatorname{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \operatorname{Pic} X^{\operatorname{al}}$.
- We can compute Aut A, the isomorphism class seems to be $F_7 \times PGL(2,7)$.
- No small rational curves: There are no lines, no conics, no twisted cubics.
- The "smallest" non-trivial curves that appear are smooth rational quartics.
- + Lattice computations with Λ predict that there are

133056

smooth rational quartics spanning Λ .

Reconstructing isolated curves from their Hodge classes

Turns out one can compute a bit more for hypersurfaces

$$\varphi: H_2(X, \mathbb{Z}) \times H^2_{\mathrm{dR}}(X/k) \to \mathbb{C} \qquad (\gamma, \omega) \longmapsto \int_{\gamma} \omega$$

Note, if $\gamma \in \operatorname{Pic} X^{\operatorname{al}}$, then $\frac{1}{2\pi i} \int_{\gamma} \omega \in k^{\operatorname{al}}$ for $\omega \in F^1 H^2_{dR}(X/k)$.

Reconstructing isolated curves from their Hodge classes

Turns out one can compute a bit more for hypersurfaces

$$\varphi: H_2(X, \mathbb{Z}) \times H^2_{\mathrm{dR}}(X/k) \to \mathbb{C} \qquad (\gamma, \omega) \longmapsto \int_{\gamma} \omega$$

Note, if $\gamma \in \operatorname{Pic} X^{\operatorname{al}}$, then $\frac{1}{2\pi i} \int_{\gamma} \omega \in k^{\operatorname{al}}$ for $\omega \in F^1 H^2_{dR}(X/k)$.

Theorem (Movasati-Sertöz)

If $\gamma = [Y] \in H_2(X, \mathbb{Z})$ for a curve $Y \subset X$ then from $\frac{1}{2\pi i} (\int_{\gamma} \omega)_{\omega \in F^1}$ one can construct an ideal I_{γ} such that $I(Y) \subsetneq I_{\gamma}$.

In favorable circumstances we expect low order equations in I_{γ} to span I(Y).

Reconstructing isolated curves from their Hodge classes

Turns out one can compute a bit more for hypersurfaces

$$\varphi: H_2(X, \mathbb{Z}) \times H^2_{\mathrm{dR}}(X/k) \to \mathbb{C} \qquad (\gamma, \omega) \longmapsto \int_{\gamma} \omega$$

Note, if $\gamma \in \operatorname{Pic} X^{\operatorname{al}}$, then $\frac{1}{2\pi i} \int_{\gamma} \omega \in k^{\operatorname{al}}$ for $\omega \in F^1 H^2_{dR}(X/k)$.

Theorem (Movasati-Sertöz)

If $\gamma = [Y] \in H_2(X, \mathbb{Z})$ for a curve $Y \subset X$ then from $\frac{1}{2\pi i} (\int_{\gamma} \omega)_{\omega \in F^1}$ one can construct an ideal I_{γ} such that $I(Y) \subsetneq I_{\gamma}$.

In favorable circumstances we expect low order equations in I_{γ} to span I(Y).

Theorem (Cifani-Pirola-Schlesinger)

For a smooth rational quartic $Y \subset X$ we have that the equation of the quadric containing Y generates $I_{[Y],2}$, i.e., $I(Y)_2 = I_{[Y],2}$.

$$\begin{aligned} X : x^4 + xyzw + y^3z + yw^3 + z^3w &= 0 \subset \mathbb{P}^3 \\ \mathsf{Pic}(X^{\mathsf{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}\,X^{\mathsf{al}} \end{aligned}$$

Goal

Reconstruct the quadrics containing some of the 133056 smooth rational quartics in *X* using the curve classes.

$$\begin{aligned} X: x^4 + xyzw + y^3z + yw^3 + z^3w &= 0 \subset \mathbb{P}^3\\ \mathsf{Pic}(X^{\mathsf{al}})_{|B} &\subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}\,X^{\mathsf{al}} \end{aligned}$$

Goal

Reconstruct the quadrics containing some of the 133056 smooth rational quartics in *X* using the curve classes.

• Fortunately, there is a small $Aut(\Lambda)$ orbit of size 336.

$$\begin{aligned} X: x^4 + xyzw + y^3z + yw^3 + z^3w &= 0 \subset \mathbb{P}^3\\ \operatorname{Pic}(X^{\operatorname{al}})_{|B} &\subseteq \Lambda \stackrel{?}{\subseteq} \operatorname{Pic} X^{\operatorname{al}} \end{aligned}$$

Goal

Reconstruct the quadrics containing some of the 133056 smooth rational quartics in *X* using the curve classes.

- Fortunately, there is a small $Aut(\Lambda)$ orbit of size 336.
- Hence, we expect an orbit of 168 quadrics each containing a pair of quartics.
- We aim reconstruct the ten (algebraic!) coefficients of these quadrics.

Goal

Reconstruct the ten coefficients of these quadrics in a Galois orbit of size 168.

• Considering all the embeddings and by clearing denominators when possible, one can reconstruct each coefficient independently.

Goal

Reconstruct the ten coefficients of these quadrics in a Galois orbit of size 168.

- Considering all the embeddings and by clearing denominators when possible, one can reconstruct each coefficient independently.
- The minimal polynomials of these elements have incredibly large height.

 $- \ 1268317331496745879603035032448157273146519836562713924560050631153969519297207668270922371313x^{147} + \cdots \\$

• Every computation must be done very selectively.

Goal

Reconstruct the ten coefficients of these quadrics in a Galois orbit of size 168.

- Considering all the embeddings and by clearing denominators when possible, one can reconstruct each coefficient independently.
- The minimal polynomials of these elements have incredibly large height.

 $x^{168} - 10014013832542203812872613924739 x^{161} + 171047690745503707515328576627906817785436888130925209472262244 x^{154} + 17104769074550370751532857662790687785478 + 1710476907455037075153884 + 171047690745503707515384 + 17104769074550370751532857662790681778543688813092520947262244 + 17104769074550370751532857662790681778543688813092520947262244 + 1710476907455037075153285766279068177854368881309252094726284 + 17104769076884 + 17104769076884 + 17104769076884 + 17104769076884 + 17104769076884 + 17104769076884 + 17104769076884 + 1710476907884 + 17104769084 + 17104769084 + 17104769084 + 17104769084 + 17104769084 + 1710476984 + 17104769084 + 17104769084 + 1710476984 + 1710476984 + 1710476984 + 1710476984 + 1710476984 + 1710476984 + 1710476984 + 17104$

 $-\ 1268317331496745879603035032448157273146519836562713924560050631153969519297207668270922371313x^{147}+\cdots$

- Every computation must be done very selectively.
- We solve the isomorphism problem between the different presentations by refining the complex embeddings and inverting a Vandermonde matrix. The abstract isomorphism problem feels hopeless otherwise.

Intersecting the quadric with X

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168$$

Goal

Show that $Q \cap X$ decomposes into two quartic curves.

• It suffices to show that the singular locus S of $Q \cap X$ consists of 10 distinct reduced points.

Intersecting the quadric with X

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168$$

Goal

Show that $Q \cap X$ decomposes into two quartic curves.

- It suffices to show that the singular locus S of $Q \cap X$ consists of 10 distinct reduced points.
- Hopeless to do this directly!
- Working over \mathbb{F}_p we find 10 distinct points. Hence, S is zero-dimensional and reduced, and deg $S \leq 10$.

Intersecting the quadric with X

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168$$

Goal

Show that $Q \cap X$ decomposes into two quartic curves.

- It suffices to show that the singular locus S of $Q \cap X$ consists of 10 distinct reduced points.
- Hopeless to do this directly!
- Working over \mathbb{F}_p we find 10 distinct points. Hence, S is zero-dimensional and reduced, and deg $S \leq 10$.
- We conclude deg S = 10 via Gotzmann regularity theorem, by checking that dim $L[x, y, z, w]_{\bullet}/l_{\bullet} = 10$ for $\bullet = 6, 7$, where V(l) = S.

$$\Lambda_{Q} := \langle [C] : C \subset \sigma(Q) \cap X, \, \sigma : L \hookrightarrow \mathbb{C} \rangle \subseteq \mathsf{Pic}(X^{\mathsf{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}X^{\mathsf{al}}$$

$$\Lambda_{Q} := \langle [C] : C \subset \sigma(Q) \cap X, \, \sigma : L \hookrightarrow \mathbb{C} \rangle \subseteq \mathsf{Pic}(X^{\mathsf{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}X^{\mathsf{al}}$$

Nonetheless, Pic X^{al} and Λ are saturated in $H_2(X, \mathbb{Z})$.

Hence, it is sufficient to show that $rk \Lambda_Q = rk \Lambda = 19$.

$$\Lambda_{Q} := \langle [C] : C \subset \sigma(Q) \cap X, \, \sigma : L \hookrightarrow \mathbb{C} \rangle \subseteq \mathsf{Pic}(X^{\mathsf{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}X^{\mathsf{al}}$$

Nonetheless, Pic X^{al} and Λ are saturated in $H_2(X, \mathbb{Z})$.

Hence, it is sufficient to show that $rk \Lambda_Q = rk \Lambda = 19$.

We can do this in two ways:

• Compute the intersections of these 336 curves with each other over \mathbb{F}_p .

$$\Lambda_{Q} := \langle [C] : C \subset \sigma(Q) \cap X, \, \sigma : L \hookrightarrow \mathbb{C} \rangle \subseteq \mathsf{Pic}(X^{\mathsf{al}})_{|B} \subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}X^{\mathsf{al}}$$

Nonetheless, Pic X^{al} and Λ are saturated in $H_2(X, \mathbb{Z})$.

Hence, it is sufficient to show that $rk \Lambda_Q = rk \Lambda = 19$.

We can do this in two ways:

- Compute the intersections of these 336 curves with each other over \mathbb{F}_p .
- Certify that these correspond to the original classes.
 Showing that there are at most 66528 distinct quadrics. Can be done over C.
 This establishes a bijection between these quadrics and the 168 pairs of quartic curve classes that they correspond to.

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L:=\mathbb{Q}(\{a_i\}_i):\mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L:=\mathbb{Q}(\{a_i\}_i):\mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

Via the identification with the original classes we have $\frac{1}{2\pi i} \left(\int_{C} \omega \right)_{\omega \in F^{1}} \in K^{21}$.

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L:=\mathbb{Q}(\{a_i\}_i):\mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

Via the identification with the original classes we have $\frac{1}{2\pi i} \left(\int_C \omega \right)_{\omega \in F^1} \in K^{21}$. These can be reconstructed in the same fashion as we reconstructed a_i . Unclear how to certify such heuristic guesses!

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

Via the identification with the original classes we have $\frac{1}{2\pi i} \left(\int_{\mathcal{C}} \omega \right)_{\omega \in F^1} \in K^{21}$.

These can be reconstructed in the same fashion as we reconstructed a_i .

Unclear how to certify such heuristic guesses!

Even if given the order $\mathcal{O} \subset K$ over which the quartics are defined over, no obvious control over denominators of $\frac{1}{2\pi i} \int_{C} \omega$.

Can one compute *K* using geometry without Gröbner basis?

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L:=\mathbb{Q}(\{a_i\}_i):\mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

The direct computation of $Gal(K/\mathbb{Q})$ looks hopeless.

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

The direct computation of $Gal(K/\mathbb{Q})$ looks hopeless. We guess that $K = F(\sqrt[14]{u})$ for a unit *u* of where *F* is defined by

$$x^{24} + x^{22} - 24x^{21} - 84x^{20} - 205x^{19} - 155x^{18} - 770x^{17} - 500x^{16} + 18916x^{15} + 36988x^{14} + 109234x^{13} + 387901x^{12} + 373961x^{11} - 18170x^{10} + 75132x^9 + 10381x^8 - 123071x^7 + 108274x^6 - 41580x^5 + 39936x^4 - 21911x^3 + 4032x^2 + 1428x + 616$$

and $Gal(F/\mathbb{Q}) = C_3 \times PGL(2,7)$ (with size 14 times smaller than Aut Pic X^{al}).

$$Q: a_0 x^2 + a_1 x y + \dots + a_9 w^2 = 0 \subset \mathbb{P}^3, \quad [L := \mathbb{Q}(\{a_i\}_i) : \mathbb{Q}] = 168$$

Goal

Compute *K* and $Gal(K/\mathbb{Q})$ acting on Λ_Q .

The direct computation of $Gal(K/\mathbb{Q})$ looks hopeless. We guess that $K = F(\sqrt[14]{u})$ for a unit *u* of where *F* is defined by

$$x^{24} + x^{22} - 24x^{21} - 84x^{20} - 205x^{19} - 155x^{18} - 770x^{17} - 500x^{16} + 18916x^{15} + 36988x^{14} + 109234x^{13} + 387901x^{12} + 373961x^{11} - 18170x^{10} + 75132x^9 + 10381x^8 - 123071x^7 + 108274x^6 - 41580x^5 + 39936x^4 - 21911x^3 + 4032x^2 + 1428x + 616$$

and $\operatorname{Gal}(F/\mathbb{Q}) = C_3 \times \operatorname{PGL}(2,7)$ (with size 14 times smaller than $\operatorname{Aut}\operatorname{Pic} X^{\operatorname{al}}$). Do we have $\operatorname{Gal}(K/\mathbb{Q}) \stackrel{?}{=} \operatorname{Aut} \Lambda$? Can we compute $\operatorname{Gal}(K/\mathbb{Q})$ by hand?

Theorem (C-Sertöz)

The quartic surface $X : x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$ has Pic $X^{al} = \Lambda$, generated by quartics over a quadratic extension of $L := \mathbb{Q}(\{a_i\}_i)$.

We are hoping to streamline this method and also figure out its applications. Hopefully, also be able handle families, e.g.,

$$X: x^4 + txyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3(\mathbb{Q}(t))$$

Do you have a challenge K3 surface for us?