Effective Computation of Hodge Cycles

Edgar Costa (MIT) November 28, 2024, University of Sydney

Slides available at **edgarcosta.org**.

Joint work with Nicholas Mascot, Jeroen Sijsling, John Voight, and Emre Can Sertöz

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Given A compute the endomorphism ring End A.

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From the equations of A determine a basis for End A and their equations in $A \times A$.

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- There are several in principle algorithms to do this over a number field. These involve, a day/night algorithm:
 - by day: search for reasonable morphisms;
 - by night: restrict your search space.

Our setup

Let C be a nice (smooth, projective, geometrically integral) curve over k of genus g given by equations. Let J be the Jacobian of C.

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But why?

- It is an interesting challenge [*citation needed*].
- If End J contains non-trivial idempotents, we can hope to decompose J into abelian varieties of smaller dimension.
- If End J is non-trivial, then this allows us to find a modular form that describes the arithmetic properties of J and C.

An analytic description of the Jacobian

Via a chosen embedding of k into \mathbb{C} and a projection into \mathbb{P}^2 , we can consider C as a Riemann surface, and

$$J_{\mathbb{C}} = H^{0}(C, \Omega_{C})^{\vee}/H_{1}(C, \mathbb{Z}) = \mathbb{C}^{g}/\Lambda,$$

where we pick an k basis for $H^0(C, \Omega_C) = k\omega_1 \oplus \ldots \oplus k\omega_g$, hence,

$$\Lambda = \left\{ \left(\int_{\gamma} \omega_1, \ldots, \int_{\gamma} \omega_g \right) \in \mathbb{C}^g : \gamma \in H_1(C, \mathbb{Z}) \right\} \cong \mathbb{Z}^{2g}.$$

In other words, J is a complex torus (plus a polarization).

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- $\cdot\,$ Using $\Lambda,$ we can hope to understand J analytically...
- and perhaps even to be able to transfer these results to the algebraic setting.

Heuristic solution

By picking a *k*-basis for $H^0(C, \Omega_C)$, we have

 $\operatorname{End}(J) = \{T \in M_g(k) \mid T\Lambda \subset \Lambda\}$

Hence, if Π is a period matrix for *C*, i.e., $\Lambda = \Pi \mathbb{Z}^{2g}$, then we are reduced to finding a \mathbb{Z} -basis of the solutions (T, R) to

$$T\Pi = \Pi R, \qquad T \in M_g(k^{al}), \quad R \in M_{2g}(\mathbb{Z}).$$

Heuristically, via lattice reduction algorithms, we can find such a \mathbb{Z} -basis.

The Galois module structure of End(J) is given via $T \in M_q(k^{al})$.

There is no obvious way to prove that our guesses are actually correct.

The reconstruction of $T \in M_g(k^{al})$ can be quite finicky, this lead to a whole library to work with $k \subset \mathbb{C}$ with fixed embedding.

Representing endomorphisms

$$\alpha_{\mathcal{C}} : \mathcal{C} \xrightarrow{AJ} J \xrightarrow{\alpha} J \dashrightarrow \mathsf{Sym}^{g}(\mathcal{C})$$
$$P \longmapsto \{Q_{1}, \dots, Q_{g}\} \Longleftrightarrow \alpha([P - P_{0}]) = \left[\sum_{i=1}^{g} Q_{i} - P_{0}\right]$$

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Given $\alpha \in M_g(k^{al})$ this divisor is a certificate of containment $\alpha \in \text{End } J^{al}$. **Theorem (C-Mascot-Sijsling-Voight)** We give an algorithm for $M_g(k^{al}) \ni \alpha \mapsto \begin{cases} true & \text{if } \alpha \in \text{End } J^{al}, \text{and a certificate } \alpha \in false & \text{if } \alpha \notin \text{End } J^{al} \end{cases}$

By interpolation via α_{C} or by locally solving a differential equation on $C \times C$.

We give an algorithm that computes $\operatorname{End} J^{\operatorname{al}}$ with a certificate $\boxed{\checkmark_{\bullet}}$.

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 $\{L_{p_1}(t) := \det(1 - t \operatorname{Frob}_p|H^1), L_{p_2}(t), \dots, L_{p_i}(t)\} \longrightarrow \text{upper bounds on } \operatorname{End} J^{al}$

- The $L_p(t)$ polynomials are as random as End J^{al} allows it.
- Two polynomials $L_p(t)$ and $L_q(t)$ suffice to obtain a sharp upperbound.
- For (p,q) in a set of positive density, but unknown apriori.

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 $\operatorname{Frob}_{p} \operatorname{mod} p^{N} \, \bigcirc_{\mathcal{H}} \, H^{1}_{\operatorname{crys}}(C, \mathbb{Z}_{p}) \longmapsto \operatorname{upper bounds on } \operatorname{End} J^{\operatorname{al}}$

- Frob_p mod p^N is a byproduct of computing $L_p(t) = \det(1 t \operatorname{Frob}_p | H^1_{MW})$.
- We check what correspondences $C \rightsquigarrow C \mod p$ lift to $C \rightsquigarrow C \mod p^N$.

Examples

- Our method works just as well for isogenies and projections.
- We have verified, decomposed and matched the 66, 158 curves over \mathbb{Q} of genus 2 in the *L*-functions and modular form database (LMFDB).
- \cdot The algorithms verify that the plane quartic

$$C: x^{4} - x^{3}y + 2x^{3}z + 2x^{2}yz + 2x^{2}z^{2} - 2xy^{2}z + 4xyz^{2}$$
$$-y^{3}z + 3y^{2}z^{2} + 2yz^{3} + z^{4} = 0$$

has complex multiplication.

• Try it:

https://github.com/edgarcosta/endomorphisms
contains friendly button-push algorithms.



"Dans la seconde partie de mon rapport, il s'agit des variétés kählériennes dites K3, ainsi nommées en l'honneur de Kummer, Kähler, Kodaira et de la belle montagne K2 au Cachemire." —André Weil (Photo credit: Waqas Anees)

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 \cdot smooth quartic surface in \mathbb{P}^3

$$X: f(x, y, z, w) = 0, \quad \deg f = 4$$

e.g. Fermat quartic surface $x^4 + y^4 + z^4 + w^4 = 0$.

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• Kummer surfaces, Kummer(A) := $\widetilde{A/\pm}$, with A an abelian surface.

Plays a similar role as End(A) for an abelian variety A

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 $\operatorname{Pic} X^{\operatorname{al}} \simeq H^{1,1}(X) \cap H^2(X,\mathbb{Z}) \subsetneq H^2(X,\mathbb{Z}) \simeq (-E_8)^2 \oplus U^3 \simeq \mathbb{Z}^{22}$ Thus, $1 \leq \operatorname{rk} \operatorname{Pic} X^{\operatorname{al}} \leq 20 = \dim H^{1,1}(X)$. A generic K3 surface has $\operatorname{rk} \operatorname{Pic} X^{\operatorname{al}} = 1$

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Picard lattice of a K3 surface

Let X be a K3 surface defined over $k \subset \mathbb{C}$. We view X also as a complex manifold. Pic $X^{al} \simeq \mathbb{Z} \langle algebraic \ curves \ in X \rangle / \langle linear \ equivalences \rangle \subset H_2(X, \mathbb{Z})$

Goal

From the equations of X, compute $\operatorname{Pic} X^{\operatorname{al}} \subset H_2(X, \mathbb{Z})$ as a $\operatorname{Gal}(k^{\operatorname{al}}/k)$ -module.

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Useful for studying the existence rational points, precisely, via a potential Brauer–Manin obstruction on *X*, as

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and Artin–Tate conjecture (proven) also gives disc Pic X^{al} modulo squares. Acessible via WeilPolynomialOfDegree2K3Surface($w^2 = f(x, y, z)$) $\operatorname{Pic} X^{\operatorname{al}} \simeq \mathbb{Z}\langle \operatorname{algebraic} \operatorname{curves} \operatorname{in} X \rangle / \langle \operatorname{linear} \operatorname{equivalences} \rangle \subset H_2(X, \mathbb{Z})$

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Over a number field there are several in principle algorithms to compute rk Pic X or even Pic X. These involve, a day/night algorithm:

- by day: find curve classes in Pic X;
- by night: restrict the ambient space for $\operatorname{Pic} X \subset H^2(X, \mathbb{Z})$.

An analytic approach

Lefschetz (1,1) theorem

A homology class $\gamma \in H_2(X, \mathbb{Z})$ is in Pic X^{al} if and only if $\int_{\gamma} \omega_X = 0$, where ω_X is the nonzero holomorphic 2-form ω_X on X, unique up to scaling.

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Hence, if $\Pi = [\int_{\gamma} \omega_X]_{\gamma \in H_2(X,\mathbb{Z})} \in \mathbb{C}^{22}$ is the period vector for ω_X , then we are reduced to finding a lattice $\Lambda \subset H_2(X,\mathbb{Z})$ of solutions

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- Heuristically, via lattice reduction algorithms, we can find $\Lambda \subset H_2(X, \mathbb{Z})$.
- There is no obvious way to prove that our guesses are actually correct.
- Nonetheless, given Π as a ball, one can compute $B \gg 0$ such that such that

 $\operatorname{Pic}(X^{\operatorname{al}})_{|B} := \mathbb{Z}\langle \gamma \in \operatorname{Pic} X^{\operatorname{al}} \mid -\gamma_{\operatorname{prim}}^2 < B \rangle \subseteq \Lambda$ (Lairez-Sertöz).

$$X: x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$$

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 - If p and q are two primes of good reduction, such that

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- No known explicit descriptions of Pic X^{al}.

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- $\cdot\,$ Lattice computations with Λ predict that there are

133056

smooth rational quartics spanning Λ .

Reconstructing isolated curves from their Hodge classes

Turns out one can compute a bit more for hypersurfaces

$$\varphi: H_2(X, \mathbb{Z}) \times H^2_{\mathrm{dR}}(X/k) \to \mathbb{C} \qquad (\gamma, \omega) \longmapsto \int_{\gamma} \omega$$

Note, if $\gamma \in \operatorname{Pic} X^{\operatorname{al}}$, then $\frac{1}{2\pi i} \int_{\gamma} \omega \in k^{\operatorname{al}}$ for $\omega \in F^1 H^2_{\operatorname{dR}}(X/k)$.

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Theorem (Movasati-Sertöz)

If $\gamma = [C] \in H_2(X, \mathbb{Z})$ for a curve $C \subset X$ then from $\frac{1}{2\pi i} (\int_{\gamma} \omega)_{\omega \in F^1}$ one can construct an ideal I_{γ} such that $I(C) \subsetneq I_{\gamma}$.

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Theorem (Cifani-Pirola-Schlesinger)

For a smooth rational quartic curve $C \subset X$ we have that the equation of the quadric surface containing C generates $I_{[C],2}$, i.e., $I(C)_2 = I_{[C],2}$.

$$\begin{aligned} X: x^4 + xyzw + y^3z + yw^3 + z^3w &= 0 \subset \mathbb{P}^3\\ \mathsf{Pic}(X^{\mathsf{al}})_{|B} &\subseteq \Lambda \stackrel{?}{\subseteq} \mathsf{Pic}\,X^{\mathsf{al}} \end{aligned}$$

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Reconstruct the quadric surfaces containing some of the 133056 smooth rational quartics in *X* using the curve classes.

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- Fortunately, there is a small $Aut(\Lambda)$ orbit of size 336 (Elkies).
- For each quartic curve $C \subset X$, we can compute

$$I_{[C],2} = \langle a_0 x^2 + \cdots \sigma(a_9) W^2 \rangle_{\mathbb{C}}$$

that defines a quadric surface Q, such that $Q \cap X = C \cup \overline{C}$. Hence, we expect an orbit of 168 quadrics each containing a pair of quartics.

• We aim reconstruct the ten (algebraic!) coefficients of these quadrics.

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- The minimal polynomials have large height about 9k characters, e.g.:
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- Every computation must be done extremely selectively!
- We are presented with same 168 degree field *L* in 9 different ways. The abstract isomorphism problem feels hopeless. 😨

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Construct
$$\mathbb{Q}(a_k) \hookrightarrow L$$
, where $L = \mathbb{Q}(a_0, \ldots, a_9) = \mathbb{Q}(a_0)$.

One approach is to compute the roots of the minimal polynomial of a_i in L. In many situations, particularly if deg $L \gg \deg \mathbb{Q}(a_k)$, it is wiser to factor the defining polynomial of L over $\mathbb{Q}(a_k)$. This is what **PARI/GP** does.

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In our case, we have all the compatible embeddings

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In practice, it is faster to iteratively refine the complex embeddings, as their height is smaller than theoretically possible: 4k vs 120k digits.

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Show that $Q \cap X$ decomposes into two quartic curves.

• It suffices to show that the singular locus S of $Q \cap X$ consists of 10 distinct reduced points.

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- We conclude deg S = 10 via Gotzmann regularity theorem, by checking that dim $L[x, y, z, w]_{\bullet}/I_{\bullet} = 10$ for $\bullet = 6, 7$, where V(I) = S.

Certifying $\operatorname{Pic} X^{\operatorname{al}} = \Lambda$

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The inclusion $\Lambda_Q \subseteq \Lambda$ is not explicit!
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To try: For a generic hyperplane $Q \cap X \cap H$ is a degree 8 reduced scheme. The number field K is the quadratic extension where we observe two orbits.

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and $\operatorname{Gal}(F/\mathbb{Q}) = C_3 \times \operatorname{PGL}(2,7)$. # $\operatorname{Gal}(F/\mathbb{Q})$ is 14 times smaller than # Aut Pic X^{al}.

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The direct computation of $Gal(K/\mathbb{Q})$ looks hopeless.

We guess that $K = F(\sqrt[14]{u})$ for a unit *u* of where *F* is defined by

 $\begin{aligned} x^{24} + x^{22} - 24x^{21} - 84x^{20} - 205x^{19} - 155x^{18} - 770x^{17} - 500x^{16} + 18916x^{15} \\ &+ 36988x^{14} + 109234x^{13} + 387901x^{12} + 373961x^{11} - 18170x^{10} + 75132x^9 + 10381x^8 \\ &- 123071x^7 + 108274x^6 - 41580x^5 + 39936x^4 - 21911x^3 + 4032x^2 + 1428x + 616 \end{aligned}$

and $\operatorname{Gal}(F/\mathbb{Q}) = C_3 \times \operatorname{PGL}(2,7)$. # $\operatorname{Gal}(F/\mathbb{Q})$ is 14 times smaller than # Aut Pic X^{al}. Can we compute $\operatorname{Gal}(K/\mathbb{Q})$? $\operatorname{Gal}(K/\mathbb{Q}) \stackrel{?}{=} \operatorname{Aut} \Lambda$? $H^1(\operatorname{Gal}(k^{\operatorname{al}}/k), \operatorname{Pic} X^{\operatorname{al}}) =$?

Theorem (C-Sertöz)

The quartic surface $X : x^4 + xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3$ has Pic $X^{al} = \Lambda$, generated by quartics over a quadratic extension of $L := \mathbb{Q}(\{a_i\}_i)$.

We are hoping to streamline this method and also figure out its applications/limitations.

Hopefully, also be able handle families, e.g.,

$$X: x^4 + \lambda xyzw + y^3z + yw^3 + z^3w = 0 \subset \mathbb{P}^3(\mathbb{Q}(\lambda))$$

Do you have a challenge K3 surface for us?